# STABILITY OF LAWVERE THEORIES

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ABSTRACT. TODO: Write the abstract.

# Contents

1. Introduction	1
1.1. Prerequisites	1
2. Lawvere theories	2
2.1. Definition	2
2.2. Models	2
2.3. Products of models	3
3. Tensor product	5
3.1. Commutativity	5
3.2. Tensor product	5
3.3. Derived theories	7
3.4. Matrix calculus	10
3.5. Center	11
4. Morita theory	12
5. Stability	13
6. Conclusion	14
References	14

# 1. INTRODUCTION

### TODO: Write the introduction

1.1. **Prerequisites.** Let  $\mathcal{C}$  be a category with products. Let I be some index set, and  $d_i \in \mathcal{C}$ , for  $i \in I$ . The projection onto the *i*-th component is denoted by  $\pi_i : \prod_j d_j \longrightarrow d_i$ . If all the  $d_i$ 's are equal to some  $d \in \mathcal{C}$ , we write  $d^I := \prod_{i \in I} d_i$ . By universal property, a collection of morphisms  $g_i : c \longrightarrow d_i$  induces a morphism

$$\langle g_i \rangle_{i \in I} : c \longrightarrow \prod_{i \in I} d_i.$$
 (1.1)

If all the  $g_i$ 's are identities, then  $\Delta_I := \langle \mathrm{id}_c \rangle_{i \in I} : c \longrightarrow c^I$  is called a *diagonal* of c. We abbreviate  $\Delta_I$  as simply  $\Delta$  if no ambiguity arise. If  $f : b \longrightarrow c$  and  $h_i : d_i \longrightarrow e_i$ , then

$$\langle h_i g_i f \rangle_i = \left(\prod_i h_i\right) \langle g_i \rangle_i f.$$
 (1.2)

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#### CÉDRIC HO THANH

In particular,  $\langle g_i \rangle_i = (\prod_i g_i) \Delta$ . A bijection  $\sigma : I \longrightarrow I$  induces an automorphism

$$\sigma \coloneqq \left\langle \pi_{\sigma(i)} \right\rangle_{i \in I} \colon d^{I} \longrightarrow d^{I}. \tag{1.3}$$

We denote by  $S_n$  the *n*-th symmetric group, which is the group of automorphisms of an *n*-element set.

### 2. Lawvere theories

2.1. **Definition.** For  $n \in \mathbb{N}$ , let [n] be the *n*-elements set  $\{1, 2, \ldots, n\}$ . If no ambiguity arises, we shall simply denote it by n = [n]. Let  $\aleph_0$  be the full subcategory of Set spanned by sets of the form *n*. It is a skeleton of  $\text{Set}_f$ , the category of finite sets.

Clearly,  $\aleph_0$  has finite coproducts, given by addition of numbers. Therefore,  $\aleph_0^{\rm op}$  has finite products, also given by addition of numbers. In particular, all objects of  $\aleph_0^{\rm op}$  are iterated products of the object 1.

A Lawvere theory (also called one-storted algebraic theory [ARV11, definition 11.3]) is a category L with finite products, equipped with a functor  $l: \aleph_0^{\text{op}} \longrightarrow \mathsf{L}$  that is an identity on objects and preserves finite products. In other words, L is a category with finite products, where all objects are iterated products of a distinguished object 1. Therefore, the functor l shall remain implicit. A morphism  $F: \mathsf{L} \longrightarrow \mathsf{K}$  between Lawvere theories is simply a functor below  $\aleph_0^{\text{op}}$ . Let  $\mathcal{L}$ aw be the category of Lawvere theories and such morphisms. It has an initial object, namely  $\aleph_0^{\text{op}}$  itself.

2.2. Models. Let C be a category with finite products. A model of L in C is a finite product preserving functor  $X : L \longrightarrow C$ . If no ambiguity arise, we write  $X^n$  instead of Xn, X instead of  $X^1 = X1$ , and if  $f \in \mathcal{L}(m, n)$ , we write  $f : X^m \longrightarrow X^n$  instead of Xf. With these notations, a model X is the datum of an object  $X \in C$  together with a morphism  $f : X^n \longrightarrow X$  for every morphism  $f \in L(n, 1)$ , satisfying the same relations as L. A morphism of models  $\alpha : X \longrightarrow Y$  is simply a natural transformation. Let L(C) be the category of models of L in C and  $U_L : L(C) \longrightarrow C$  be the forgetful functor. It maps a model X to  $X = X^1$ , and  $\alpha : X \longrightarrow Y$  to  $\alpha = \alpha_1$ .

**Lemma 2.1.** If  $\alpha : X \longrightarrow Y$  is a morphism of models, then  $\alpha_n : X^n \longrightarrow Y^n$  is simply  $(\alpha_1)^n$ .

*Proof.* Consider the projection  $\pi_i : n \longrightarrow 1$  in L. Since  $\alpha$  is a natural transformation, the following square commutes:

Since this holds for all  $1 \le i \le n$ , we conclude that  $\alpha_n = (\alpha_1)^n$ .

**Lemma 2.3.** Let  $X, Y \in L(\mathbb{C})$  and  $\beta : X \longrightarrow Y$  be a morphism in  $\mathbb{C}$ . Then  $\beta$  extends as a morphism of models  $X \longrightarrow Y$  if and only if for all  $n \in \mathbb{N}$  and  $f \in L(n, 1)$ , the

following square commutes

*Proof.* Necessity is clear. We now prove that the condition is sufficient. Note that a morphism  $g: n \longrightarrow m$  in L can be decomposed as  $g = \langle \pi_i g \rangle_{1 \le i \le m}$ , and that each  $\pi_i g$  is a morphism  $n \longrightarrow 1$ . By assumption,  $\pi_i g \beta^n = \beta \pi_i g$ , thus

$$g\beta^{n} = \langle \pi_{i}g\beta^{n} \rangle_{1 \le i \le m} = \langle \beta\pi_{i}g \rangle_{1 \le i \le m} = \beta g.$$

$$(2.5)$$

Thanks to lemmas 2.1 and 2.3, we can write the *n*-th component of  $\alpha : X \longrightarrow Y$ simply as  $\alpha^n : X^n \longrightarrow Y^n$ . Consider the case  $\mathcal{C} = \text{Set}$ , and let  $n \in \mathbb{N}$ . The *n*-th representable model of L [ARV11, remark 1.12] is  $L_n := L(n, -) : L \longrightarrow \text{Set}$ .

**Lemma 2.6** (Yoneda lemma). For  $X \in L(Set)$  there is an natural isomorphism (of sets)

$$X^n \cong \mathsf{L}(\mathsf{Set})(L_n, X). \tag{2.7}$$

More precisely, it maps an element  $x \in X^n$  to the unique morphism  $\tilde{x} : L_n \longrightarrow X$ such that  $\tilde{x}^n(\operatorname{id}_n) = x$ .

2.3. Products of models. By [AR94, theorem 1.46 and corollary 1.52], L(Set) is a finitely locally presentable category, and in particular, it has all limits and colimits. The same result cannot be expected to hold for a general ground category C with finite products, but as we will see in lemma 2.11, L(C) still has finite products.

For integers m, n > 0, let  $\tau_{n;m} \in S_{mn}$  be the *shuffle permutation*, i.e. the permutation such that for all  $1 \le i \le m$  and  $1 \le j \le n$ , we have

$$\tau_{n;m}((j-1)n+i) = (i-1)m+j.$$
(2.8)

In other words,  $\tau_{n;m}$  "rearranges m tuples of n elements into n tuples of m elements". If n or m is 0, then by convention, let  $\tau_{n;m}$  be the identity on the empty set. We consider  $\tau_{n;m}$  as a morphism  $mn \longrightarrow mn$  of  $\aleph_0^{\text{op}}$ , and therefore as a morphism in any Lawvere theory.

If  $c_1, \ldots, c_n \in \mathbb{C}$  and  $m \in \mathbb{N}$ , then there is an isomorphism

$$\tau = \tau_{c_1,\dots,c_n;m} : \left(\prod_i c_i\right)^m \longrightarrow \prod_i c_i^m \tag{2.9}$$

induced by  $\tau_{n;m}$ , which again, rearranges m tuples of n elements into n tuples of m elements. Explicitly, if  $p_k : (\prod_i c_i)^m \longrightarrow c_l$  is the projection into the k-th component, where  $1 \le k \le mn$ ,  $1 \le l \le n$ , and  $l \equiv k \mod n$ , then  $\tau$  is the universal morphism

$$\tau = \left\langle p_{\tau_{n;m}(k)} \right\rangle_{1 \le k \le mn}. \tag{2.10}$$

If all the  $c_i$ 's are equal to some object c, then we write  $\tau_{c:n:m}$  instead of  $\tau_{c...,c:m}$ .

**Lemma 2.11** ([ARV11, proposition 1.21]). *The category* L(C) *has finite products, given as follows:* 

 if the terminal object of C is 1, then the terminal model is the constant functor at 1;

(2) if 
$$X, Y \in L(\mathbb{C})$$
, then  $X \times Y$  is the functor  
 $X \times Y : L \longrightarrow \mathbb{C}$   
 $n \longmapsto (X \times Y)^n$   $n \in \mathbb{N}$   
 $f = e^{-1}$   $(f \oplus f) = e^{-1}$ 

$$f \longmapsto \tau_{X,Y;q}^{-1}(f \times f)\tau_{X,Y;p} \qquad \qquad f \in \mathsf{L}(p,q).$$

In details,  $(X \times Y)f$  is the composite

$$(X \times Y)^p \xrightarrow{\tau_{X,Y;p}} X^p \times Y^p \xrightarrow{f \times f} X^q \times Y^q \xrightarrow{\tau_{X,Y;q}^{-1}} (X \times Y)^q.$$
(2.12)

*Proof.* The model described in point (1) is clearly terminal. Consider  $X \times Y$  as defined in point (2), and define  $\pi_{X,n}$  as the composite

$$(X \times Y)^n \xrightarrow{\tau_{X,Y;n}} X^n \times Y^n \xrightarrow{\pi_{X^n}} X^n.$$
(2.13)

Take  $f \in L(p,q)$ , and consider

$$(X \times Y)^{p} \xrightarrow{\tau_{X,Y;p}} X^{p} \times Y^{p} \xrightarrow{\pi_{X}p} X^{p}$$

$$\downarrow (X \times Y)^{f} \qquad \downarrow f \times f \qquad \downarrow f$$

$$(X \times Y)^{q} \xrightarrow{\tau_{X,Y;q}} X^{q} \times Y^{q} \xrightarrow{\pi_{X}q} X^{q}.$$

$$(2.14)$$

Both inner squares commute by definition, thus so does the outer one. In other words,  $\pi_{X,p}$ 's jointly define a natural transformation  $\pi_X : X \times Y \longrightarrow X$ . Similarly, we define  $\pi_Y : X \times Y \longrightarrow Y$ . We now check that  $X \xleftarrow{\pi_X} X \times Y \xrightarrow{\pi_Y}$  is a limit cone. For  $X \xleftarrow{\alpha} Z \xrightarrow{\beta} Y$  another cone, let  $\gamma_n$  be the following composite

$$Z^{n} \xrightarrow{(\alpha_{n},\beta_{n})} X^{n} \times Y^{n} \xrightarrow{\tau_{X,Y;n}^{-1}} (X \times Y)^{n}.$$
(2.15)

Akin to (2.14), it is easy to check that the  $\gamma_n$ 's jointly define a morphism  $Z \longrightarrow X \times Y$ , and that furthermore,  $\alpha = \pi_X \gamma$  and  $\beta = \pi_Y \gamma$ . If  $\gamma'$  is another such morphism, then necessarily

$$\tau_{X,Y;2}\gamma'_n = \langle \alpha_n, \beta_n \rangle = \tau_{X,Y;2}\gamma, \qquad (2.16)$$

and since  $\tau_{X,Y;2}$  is an isomorphism,  $\gamma'_n = \gamma_n$ . Finally,  $X \times Y$  defined above is the product of X and Y.

**Lemma 2.17.** If  $X \in L(\mathcal{C})$ ,  $g \in L(n, 1)$ , and  $m \in \mathbb{N}$ , then

$$X^{m}g = (Xg)^{m}\tau_{X;m;n} = X(g^{m}\tau_{m;n}).$$
(2.18)

*Proof.* The second equality follows from the fact that X preserves finite products. For the first one, we proceed by induction. The cases m = 0, 1 hold trivially, and m = 2 holds by definition of the product in L(C). Thus, assume that  $m \ge 3$ . We

$$\begin{split} X^{m}g &= (X \times X^{m-1})g \\ &= \tau_{X,X^{m-1};1}^{-1} (Xg \times X^{m-1}g)\tau_{X,X^{m-1};1} \\ &= (Xg \times X^{m-1}g)\tau_{X,X^{m-1};1} \\ &= (Xg \times (Xg)^{m-1}\tau_{X;m-1;n})\tau_{X,X^{m-1};1} \\ &= (Xg)^{m} (\operatorname{id}_{X} \times \tau_{X;m-1;n})\tau_{X,X^{m-1};1} \\ &= (Xg)^{m} \tau_{X;m;n}. \end{split}$$

#### 3. Tensor product

3.1. Commutativity. Let  $K, L \in \mathcal{L}$ aw, and  $\mathcal{C}$  be a category with finite products. As we saw in lemma 2.11,  $L(\mathcal{C})$  has finite products, and therefore, we may consider the category of K-models in  $L(\mathcal{C})$ , i.e.  $K(L(\mathcal{C}))$ . As we will see in this section, the latter can be seen as the category of models in  $\mathcal{C}$ .

Let  $f \in L(n, 1)$  and  $g \in L(m, 1)$ . We say that f commutes with g, denoted by  $f \boxtimes g$ , if the following equality is satisfied:

$$fg^n = gf^m \tau_{m;n}. \tag{3.1}$$

More generally, if  $c \in \mathbb{C}$ ,  $f \in \mathbb{C}(c^n, 1)$  and  $g \in \mathbb{C}(c^m, 1)$ , then we say that f commutes with g if  $fg^n = gf^m \tau_{c;m;n}$ , i.e. if the following diagram commutes:

Note that since  $\tau_{m;n}^{-1} = \tau_{n;m}$ , the commutativity relation  $\boxtimes$  is symmetric.

To simplify notations, if  $A, B \subseteq L/1$  are sets of morphisms with codomain 1, then we write  $A \boxtimes B$  to signify that every morphism in A commutes with every morphism of B. Let  $L|_A$  be the smallest subtheory of L containing A. In particular,  $L|_A/1$  is the smallest subset of L/1 such that

- (1)  $A \subseteq L|_A/1;$
- (2) the projection  $\pi_i : n \longrightarrow 1$  is in  $L|_A/1$ , for all  $n \in \mathbb{N}$  and  $1 \le i \le n$ ;
- (3) if  $g, f_1, \ldots, f_n \in L|_A/1$ , then  $g \prod_i f_i \in L|_A/1$  and  $g \langle f_i \rangle \in L|_A/1$ .

If  $L = L|_A$ , then we say that A generates L. Trivially L/1 generates L.

**Lemma 3.3.** Let  $A, B \subseteq L/1$ . If  $A \boxtimes B$ , then  $(L|_A/1) \boxtimes (L|_B/1)$ .

*Proof.* By symmetry of  $\boxtimes$ , it is enough to show that  $A \boxtimes (L|_B/1)$ . The rest is routine verifications.

3.2. Tensor product. For  $L, K \in \mathcal{L}$ aw, consider the theory  $K \coprod_{\aleph_0^{\mathrm{op}}} L$  given by the categorical pushout

$$\begin{array}{c} \aleph_0^{\text{op}} & \stackrel{k}{\longrightarrow} & \mathsf{K} \\ \iota & \downarrow & \downarrow \\ \mathsf{L} & \stackrel{\Gamma}{\longrightarrow} & \mathsf{K} \coprod_{\aleph_0^{\text{op}}} \mathsf{L}. \end{array}$$
(3.4)

have

#### CÉDRIC HO THANH

It contains all the morphisms of K and L, but is quotiented in such a way that only one copy of  $\aleph_0^{\text{op}}$  is present. In particular,  $\mathsf{K} \coprod_{\aleph^{\text{op}}} \mathsf{L}$  is still a Lawvere theory. Let  $K \otimes L$ , the tensor product (also called Kronecker product) of K and L [Fre66] be the following quotient:

$$\mathsf{K} \otimes \mathsf{L} \coloneqq \frac{\mathsf{K} \coprod_{\aleph_0^{\mathrm{op}}} \mathsf{L}}{fg^n \sim gf^m \tau_{m;n}, \ f \in \mathsf{K}(n,1), g \in \mathsf{L}(m,1)}.$$
(3.5)

Denote by  $i_{\mathsf{K}}:\mathsf{K}\longrightarrow\mathsf{K}\otimes\mathsf{L}$  the natural map (which is not necessarily faithful!), and likewise for L. Since  $\tau_{m;n}^{-1} = \tau_{n;m}$ , we have  $fg^n = gf^m \tau_{m;n}$  if and only if  $gf^m \sim fg^n \tau_{n;m}$ , which implies that the tensor product  $\otimes$  is commutative.

**Proposition 3.6.** The initial theory  $\aleph_0^{\text{op}}$  is a neutral element for the tensor product. *Proof.* Take  $\mathsf{K} \in \mathcal{L}$ aw. Then  $\mathsf{K} \coprod_{\aleph_0^{\mathrm{op}}} \aleph_0^{\mathrm{op}} = \mathsf{K}$ , thus

$$\begin{split} \mathsf{K} \otimes \aleph_0^{\mathrm{op}} &= \frac{\mathsf{K}}{fg^n \sim gf^m \tau_{m;n}, \ f \in \mathsf{K}(n,1), g \in \aleph_0^{\mathrm{op}}(m,1)} \\ &= \frac{\mathsf{K}}{f\pi_i^n \sim \pi_i f^m \tau_{m;n}, \ f \in \mathsf{K}(n,1), \pi_i : m \longrightarrow 1, 1 \leq i \leq m} \end{split}$$

However, the relation  $f\pi_i^n = \pi_i f^m \tau_{m;n}$  is already satisfied in K (and indeed, in any Lawvere theory), thus  $\mathsf{K} \otimes \aleph_0^{\mathrm{op}} = \aleph_0^{\mathrm{op}} \otimes \mathsf{K} = \mathsf{K}$ .

# Lemma 3.7. Let $X \in K(L(\mathcal{C}))$ .

- (1) For d ∈ hom ℵ<sub>0</sub><sup>op</sup> we have (X1)d = (Xd)<sub>1</sub>.
  (2) For f ∈ K/1 and g ∈ L/1 we have (X1)g ⊠ (Xf)<sub>1</sub>.
- Proof. (1) Since X and X1 are models of Lawyere theories, they preserve finite products.
  - (2) By naturality of Xf, the following diagram commutes

$$(Xm)n \xrightarrow{(Xf)_n} (X1)n$$

$$(Xm)g \downarrow \qquad \qquad \downarrow (X1)g$$

$$(Xm)1 \xrightarrow{(Xf)_1} (X1)1.$$
(3.8)

By lemma 2.11,  $(Xm)g = (X1)^m g = ((X1)g)^m \tau_{X1;m;n}$ , thus

$$(Xf)_1((X1)g)^m \tau_{(X1)1;m;n} = (X1)g(Xf)_n = ((X1)g)(Xf)_1^n, \qquad (3.9)$$

which is the desired relation.

**Theorem 3.10.** There is an equivalence  $K(L(\mathcal{C})) \simeq K \otimes L(\mathcal{C})$ .

*Proof.* Informally, the central observation is that the relations  $fg^n = gf^m \tau_{m:n}$  ranging over q encodes the fact that f induces a natural transformation between Lmodels in  $\mathcal{C}$ , i.e. a morphism in  $L(\mathcal{C})$ .

Let  $X \in \mathsf{K}(\mathsf{L}(\mathcal{C}))$ . By lemma 3.7, the following functor is a well-defined model of  $\mathsf{K}\otimes\mathsf{L}$ :

$$\begin{aligned} X^{\flat} &: \mathsf{K} \otimes \mathsf{L} \longrightarrow \mathfrak{C} \\ & n \longmapsto (X1)n & n \in \mathbb{N} \\ & f \longmapsto (Xf)_1 & f \in \hom \mathsf{K} \\ & g \longmapsto (X1)g & g \in \hom \mathsf{L}. \end{aligned}$$

check

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This defines a functor  $(-)^{\flat} : \mathsf{K}(\mathsf{L}(\mathcal{C})) \longrightarrow \mathsf{K} \otimes \mathsf{L}(\mathcal{C})$  mapping a morphism  $\alpha$  to  $\alpha^{\flat} := \alpha_1$ . Conversely, let  $Y \in \mathsf{K} \otimes \mathsf{L}(\mathcal{C})$ . For  $f \in \mathsf{L}(m, m')$  and  $g \in \mathsf{L}(n, 1)$ , the following square commutes

$$\begin{array}{cccc}
Y^{mn} & \xrightarrow{(Yf)^n} & Y^{m'n} \\
Y^mg & & & \downarrow & & \downarrow & & \\
Y^m & \xrightarrow{Yf} & & & Y^{m'}g \\
\end{array} \tag{3.11}$$

since in  $\mathsf{K} \otimes \mathsf{L}$  we have  $\pi_i f \boxtimes g$  for all  $1 \leq i \leq m$ . Therefore, by lemma 2.3, the  $(Yf)^n$ 's jointly define a morphism  $(Yf)^{\bullet} : Y^m i_{\mathsf{K}} \longrightarrow Y^{m'} i_{\mathsf{K}}$ . This enables the following definition

$$Y^{\sharp}: \mathsf{K} \longrightarrow \mathsf{L}(\mathfrak{C})$$
$$n \longmapsto Y^{n}i_{\mathsf{K}} \qquad n \in \mathbb{N}$$
$$f \longmapsto (Yf)^{\bullet} \qquad \qquad f \in \hom \mathsf{K},$$

from which we derive a functor  $(-)^{\sharp} : \mathsf{K} \otimes \mathsf{L}(\mathcal{C}) \longrightarrow \mathsf{K}(\mathsf{L}(\mathcal{C}))$  mapping a morphism  $\beta$  to  $(\beta_1)^{\bullet}$ . Finally, it is routine verification to show that  $(-)^{\flat}$  and  $(-)^{\sharp}$  are mutually inverse equivalences of categories.

**Corollary 3.12.** We have  $K(L(\mathcal{C})) \simeq L(K(\mathcal{C}))$ .

### 3.3. Derived theories.

Let Mon be the theory of monoids. It is generated by a constant  $0: 0 \rightarrow 1$ , a for that binary operation  $\lambda = \lambda_2: 2 \rightarrow 1$ , and the following relations:

$$\lambda(\mathrm{id}_1 \times 0) = \mathrm{id}_1 = \lambda(0 \times \mathrm{id}_1), \qquad \lambda(\lambda \times \mathrm{id}_1) = \lambda(\mathrm{id}_1 \times \lambda), \tag{3.13}$$

standing for unitality and associativity, respectively. Define the *n*-ary multiplication  $\lambda_n$  as follows:  $\lambda_0 \coloneqq 0$ ,  $\lambda_1 \coloneqq \operatorname{id}_1$ , and for  $n \ge 2$ ,  $\lambda_{n+1} \coloneqq \lambda(\lambda_n \times \operatorname{id}_1)$ . Let **cMon** be the theory of commutative monoids, generated like **Mon** but with the following additional commutativity axiom:

$$\lambda = \lambda(1\,2).\tag{3.14}$$

Finally, the theory Ab of abelian groups extends cMon with a morphism  $i: 1 \longrightarrow 1$ and the following invertibility axiom

$$\lambda(i \times \mathrm{id}_1) = 0 = \lambda(\mathrm{id}_1 \times i). \tag{3.15}$$

**Proposition 3.16** (Eckmann–Hilton argument). Let K be a Lawvere theory,  $n \ge 2$ ,  $f, g: n \longrightarrow 1$ , and  $0: 0 \longrightarrow 1$ . Assume that 0 is a neutral element for f and g, i.e.  $f(0^{i-1} \times id_1 \times 0^{n-i}) = id_1$ , for all  $1 \le i \le n$ , and likewise for g. If  $f \boxtimes g$ , then f = g and for all  $\sigma \in S_n$ ,  $f = f\sigma$ .

*Proof.* Take  $\sigma \in S_n$ , and define

$$x_{i,j} \coloneqq \begin{cases} \operatorname{id}_1 & \text{if } i = \sigma(j) \\ 0 & \text{otherwise.} \end{cases}$$
(3.17)

find better term

We have

$$\begin{aligned} f &= f \prod_{i} x_{i,\sigma(i)} \\ &= f \prod_{i} g \prod_{j} x_{i,j} \\ &= f g^{n} \prod_{i,j} x_{i,j} \\ &= g f^{n} \tau_{n;n} \prod_{i,j} x_{i,j} \\ &= g f^{n} \left( \prod_{i,j} x_{j,i} \right) \tau_{n;n} \\ &= g \left( \prod_{i} f \prod_{j} x_{j,i} \right) \tau_{n;n} \\ &= g \left( \prod_{i} x_{\sigma(i),i} \right) \tau_{n;n} \\ &= g \left( \prod_{i} x_{\sigma(i),i} \right) \tau_{n;n} \end{aligned}$$
 since 0 neutral for f  
$$&= g \left( \prod_{i} x_{\sigma(i),i} \right) \tau_{n;n} \\ &= g \sigma \end{aligned}$$

since  $\tau_{n;n}((i-1)n + \sigma(i)) = (\sigma(i) - 1)n + i$ . Taking  $\sigma = \mathrm{id}_n$  we find f = g. Then, for a general  $\sigma$ , we have  $f = g\sigma = f\sigma$ .

**Proposition 3.18** (Classical Eckmann–Hilton argument [EH62]). For all  $n \ge 2$ , we have cMon = Mon<sup> $\otimes n$ </sup>. In particular, cMon = cMon $\otimes$  cMon.

*Proof.* In Mon  $\otimes$  Mon, denote by 0 and  $\lambda$  the operation from the left instance of Mon, while 0' and  $\lambda'$  are those from the right instance. Since  $0 \boxtimes 0'$ , we immediately conclude that 0 = 0'. Then, 0 is neutral for both  $\lambda$  and  $\lambda'$ , and by definition,  $\lambda \boxtimes \lambda'$ . Thus, by proposition 3.16,  $\lambda = \lambda'$  and  $\lambda$  is commutative. In other words, cMon = Mon  $\otimes$  Mon. Similarly, cMon = cMon  $\otimes$  Mon.

Let  $K, L \in \mathcal{L}aw$ . We say that K is a *simple* L-*theory* if it can be decomposed as  $K = K' \otimes L$ , for some  $K' \in \mathcal{L}aw$ . It is an L-*theory* if it is the quotient of a simple L-theory.

Let K be an Mon-theory, i.e. a quotient of a theory of the form  $K' \otimes Mon$ . The set K(1,1) has a semiring structure, where multiplication is given by composition, and where addition is defined as

$$x + y \coloneqq \lambda \langle x, y \rangle \tag{3.19}$$

for  $x, y: 1 \longrightarrow 1$ . We denote this semiring by  $\varepsilon_1 \mathsf{K}$ . It is easy to check that the *n*-ary sum of  $x_1, \ldots, x_n: 1 \longrightarrow 1$  is

$$\sum_{i} x_{i} = \lambda_{n} \langle x_{i} \rangle_{i} \,. \tag{3.20}$$

**Proposition 3.21.** Ab-theories are exactly of the form  $Mod_R$ , where R is a commutative ring.

*Proof.* If K is an Ab-theory, then it is easy to see that  $K = Mod_R$ , where  $R \coloneqq \varepsilon_1 K$ .

**Lemma 3.22.** Let  $A, B \subseteq K/1$  be such that (1)  $A \cup B$  generates K;

The terminology isn't great, as it suggests enrichement.

probably not the

best place for it

probably similar results for

Mon- and cMon-

theories.

(2)  $A \boxtimes B$ .

Then K is a quotient of  $K|_A \otimes K|_B$ .

*Proof.* Since  $A \boxtimes B$  in K, the natural functors  $\mathsf{K}|_A \longrightarrow \mathsf{K}$  and  $\mathsf{K}|_B \longrightarrow \mathsf{K}$  extend to a functor  $F : \mathsf{K}|_A \otimes \mathsf{K}|_B \longrightarrow \mathsf{K}$ . Since  $A \cup B$  generates  $\mathsf{K}$ , F is full.  $\Box$ 

**Proposition 3.23.** Let  $K, L \in \mathcal{L}aw$ , where K is commutative. Then L is a K-theory if and only if it there exists a morphism  $F : K \longrightarrow L$  such that  $F(K/1) \boxtimes L/1$ .

*Proof.* (1) ( ⇒ ) Assume that L is a K-theory, i.e. obtained as a quotient of a simple K-theory, say L' = L'' ⊗ K. Let F be the composite

$$\mathsf{K} \longrightarrow \mathsf{L}'' \otimes \mathsf{K} \longrightarrow \mathsf{L}. \tag{3.24}$$

Let  $A := \mathsf{K}/1$  and  $B := \mathsf{L}''/1$ . Then by definition of the tensor product, in  $\mathsf{L}'$ , we have  $A \boxtimes B$ . Furthermore, by commutativity, in  $\mathsf{K}$  we have  $A \boxtimes A$ . Thus, in  $\mathsf{L}'$ , we have  $A \boxtimes A \cup B$ . Since  $\mathsf{L}'$  is generated by  $A \cup B$ , it follows from lemma 3.3 that  $A \boxtimes \mathsf{L}'/1$ . Finally,  $F(\mathsf{K}/1) = F(A) \boxtimes F(\mathsf{L}'/1) = \mathsf{L}/1$ .

(2) (  $\Leftarrow$  ) Note that  $A \coloneqq F(K/1)$  and  $B \coloneqq L/1$  satisfy the requirements of lemma 3.22.

**Corollary 3.25.** A theory K is a cMon-theory if and only if it contains cMon and  $\{0, \lambda\} \boxtimes K/1$ .

*Proof.* By the Eckmann–Hilton argument (proposition 3.18), cMon is commutative. Apply proposition 3.23.  $\Box$ 

Assume that K is a Mon-theory, and let  $f \in K(n, 1)$ . For  $1 \le i \le n$ , the *i*-th axis  $f^{[i]}$  of f [BV79] is defined as

$$f^{[i]} \coloneqq f(0^{i-1} \times \mathrm{id}_1 \times 0^{n-i}). \tag{3.26}$$

**Proposition 3.27.** We have  $f = \lambda_n \prod_{i=1}^n f^{[i]}$ . Further, the  $f^{[i]}$ 's are unique for this property.

*Proof.* We have

$$f = f \prod_{i=1}^{n} \lambda_n (0^{i-1} \times \mathrm{id}_1 \times 0^{n-i})$$
  

$$= f \lambda_n^n \prod_{i=1}^{n} (0^{i-1} \times \mathrm{id}_1 \times 0^{n-i})$$
  

$$= \lambda_n f^n \tau_{n;n} \prod_{i=1}^{n} (0^{i-1} \times \mathrm{id}_1 \times 0^{n-i})$$
  

$$= \lambda_n f^n \prod_{i=1}^{n} (0^{i-1} \times \mathrm{id}_1 \times 0^{n-i})$$
  

$$= \lambda_n \prod_{i=1}^{n} f(0^{i-1} \times \mathrm{id}_1 \times 0^{n-i})$$
  

$$= \lambda_n \prod_{i=1}^{n} f^{[i]}.$$

To be precise, we need to invoke proposition 5.2

If 
$$f = \lambda_n \prod_{i=1}^n f_i$$
, for some  $f_1, \dots, f_n : 1 \longrightarrow 1$ , then  

$$f^{[i]} = f(0^{i-1} \times \operatorname{id}_1 \times 0^{n-i})$$

$$= \left(\lambda_n \prod_{j=1}^n f_j\right) (0^{i-1} \times \operatorname{id}_1 \times 0^{n-i})$$

$$= \lambda_n \left( \left( \prod_{j=1}^{i-1} f_j 0 \right) \times f_i \times \left( \prod_{j=i+1}^n f_j 0 \right) \right)$$

$$= \lambda_n (0^{i-1} \times f_i \times 0^{n-i}) \qquad \text{since } f_j \boxtimes 0$$

$$= f_i.$$

3.4. Matrix calculus. Let K be a Mon-theory. Let  $f = (f_{i,j})_{1 \le i \le m, 1 \le j \le n}$  be a  $m \times n$  matrix whose entries are elements of K(1, 1). It induces a morphism  $n \longrightarrow m$  (note the order), as

$$\left(\lambda_n \prod_{j=1}^n f_{i,j} : n \longrightarrow 1\right)_{1 \le i \le m}$$
(3.28)

**Lemma 3.29.** Conversely, a morphism  $g: n \longrightarrow m$  decomposes as  $(g_{i,j})_{1 \le i \le m, 1 \le j \le n}$ , where

$$g_{i,j} \coloneqq (\pi_i g)^{\lfloor j \rfloor}. \tag{3.30}$$

Therefore, there is a bijection  $M_{m \times n}(\mathsf{K}(1,1)) \cong \mathsf{K}(n,m)$ .

*Proof.* We have,  $g = \langle \pi_i g \rangle_{1 \le i \le m}$ , and by proposition 3.27,  $\pi_i g = \lambda_n \prod_{j=1}^n (\pi_i g)^{[j]}$ . The bijection follows from the fact that a morphism  $g: n \longrightarrow m$  is uniquely determined by its projections  $\pi_i g$ , which are themselves uniquely determined by their axes, see proposition 3.27.

In particular, if f is a morphism  $m \longrightarrow 1$ , then its matrix has a single row whose entries are the axes of f.

**Lemma 3.31.** If  $f = (f_{i,j})_{1 \le i \le m, 1 \le j \le n} : n \longrightarrow m$  and  $g = (g_{j,k})_{1 \le j \le n, 1 \le k \le p} : p \longrightarrow n$ , then the composite  $fg : p \longrightarrow m$  is given by the expected matrix product

$$(fg)_{i,k} = \sum_{j} f_{i,j}g_{j,k}.$$
 (3.32)

Proof. We have

$$(fg)_{i,k} = (\pi_i fg)^{[k]}$$
by lemma 3.29  

$$= \pi_i fg(0^{k-1} \times \mathrm{id}_1 \times 0^{p-k})$$
by definition  

$$= \pi_i f \langle \pi_j g \rangle_j (0^{k-1} \times \mathrm{id}_1 \times 0^{p-k})$$
  

$$= \pi_i f \langle \pi_j g(0^{k-1} \times \mathrm{id}_1 \times 0^{p-k}) \rangle_j$$
  

$$= \pi_i f \langle g_{j,k} \rangle_j$$
  

$$= \lambda_n \left( \prod_j f_{i,j} \right) \langle g_{j,k} \rangle_j$$
  

$$= \lambda_n \langle f_{i,j} g_{j,k} \rangle_j$$
  

$$= \sum_j f_{i,j} g_{j,k}.$$

**Lemma 3.33.** Let  $\sigma \in S_n$  be a permutation. The matrix of the induced morphism  $\sigma: n \longrightarrow n$  is given by

$$(\sigma_{i,j})_{1 \le i,j \le n} \coloneqq \begin{cases} \operatorname{id}_1 & \text{if } i = \sigma(j) \\ 0 & \text{otherwise.} \end{cases}$$
(3.34)

*Proof.* For the time being, let  $(s_{i,j})_{1 \le i,j \le n}$  be the matrix defined above, and s be the morphism  $n \longrightarrow n$  it induces. By definition,

$$s = \left(\lambda_n \prod_j s_{i,j}\right)_i = \left\langle \pi_{\sigma(i)} \right\rangle_i = \sigma \tag{3.35}$$

The matrix of  $id_n \in S_n$  is thus the identity matrix, also denoted by  $I_n$ :

$$I_n \coloneqq \begin{pmatrix} \operatorname{id}_1 & \cdots & 0\\ \vdots & \ddots & \vdots\\ 0 & \cdots & \operatorname{id}_1 \end{pmatrix}.$$
(3.36)

- **Lemma 3.37.** (1) For  $f: m \to 1$ ,  $\sigma \in S_m$ , and  $1 \le i \le m$ , we have  $(f\sigma)^{[i]} = f^{[\sigma^{-1}(i)]}$ .
  - (2) For  $f: m \to 1$ ,  $g: n \to 1$ ,  $1 \leq i \leq m$ , and  $1 \leq j \leq n$ , we have  $(fg^m)^{[(i-1)n+j]} = f^{[i]}g^{[j]}$ . In particular,  $f \boxtimes g$  if and only if all the axes of f commute (multiplicatively) with all the axes of g.
- *Proof.* (1) We have

$$f\sigma = (f^{[1]} \cdots f^{[m]})\sigma = (f^{[\sigma^{-1}(1)]} \cdots f^{[\sigma^{-1}(m)]}).$$
(3.38)

(2) We have

$$\begin{split} fg^{m} &= \begin{pmatrix} f^{[1]} & \cdots & f^{[m]} \end{pmatrix} \begin{pmatrix} g & \cdots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \cdots & g \end{pmatrix} \\ &= \begin{pmatrix} f^{[1]} & \cdots & f^{[m]} \end{pmatrix} \begin{pmatrix} g^{[1]} & \cdots & g^{[n]} & 0 & \cdots & \cdots & 0 \\ \vdots & & \ddots & & \vdots \\ 0 & \cdots & \cdots & 0 & g^{[1]} & \cdots & g^{[n]} \end{pmatrix} \\ &= \begin{pmatrix} f^{[1]}g^{[1]} & f^{[1]}g^{[2]} & \cdots & f^{[1]}g^{[n]} & f^{[2]}g^{[1]} & \cdots & f^{[m]}g^{[n]} \end{pmatrix}. \end{split}$$

3.5. Center. Let  $\mathsf{K} \in \mathcal{L}$ aw. Its *center*  $Z(\mathsf{K})$  is the subtheory of  $\mathsf{K}$  generated by those morphisms  $f \in \mathsf{K}/1$  such that  $f \boxtimes \mathsf{K}/1$ . We say that  $\mathsf{K}$  is *commutative* if  $Z(\mathsf{K}) = \mathsf{K}$ . In other words  $f \in Z(\mathsf{K})(m,n)$  if and only if  $f : X^m \longrightarrow X^n$  is a morphism of models, for all  $X \in \mathsf{K}(\mathsf{Set})$ . By lemma 3.3,  $\mathsf{K}$  is commutative if and only if  $\mathsf{K}/1 \boxtimes \mathsf{K}/1$ .

**Lemma 3.39.** Let  $\mathsf{K}, \mathsf{L} \in \mathcal{L}$ aw. Then  $Z(\mathsf{K}) \otimes Z(\mathsf{L}) \subseteq Z(\mathsf{K} \otimes \mathsf{L})$ .

*Proof.* Let  $f \in Z(K)$ . Then by definition,  $f \boxtimes K/1$  in K, thus the same holds in  $K \otimes L$ . Further,  $f \boxtimes L/1$  in  $K \otimes L$ . Finally,  $f \in Z(K \otimes L)$ . The same reasoning applies if  $f \in Z(L)$ .

**Proposition 3.40.** Let  $\operatorname{id} = \operatorname{id}_{\mathsf{K}(\operatorname{Set})}$  be the identity functor of  $\mathsf{K}(\operatorname{Set})$ . There is a bijection  $Z(\mathsf{K})(m,n) \longrightarrow [\operatorname{id}^m, \operatorname{id}^n]$  which induces an equivalence  $Z(\mathsf{K}) \longrightarrow \mathfrak{I}$ , where  $\mathfrak{I}$  is the full subcategory of  $[\mathsf{K}(\operatorname{Set}), \mathsf{K}(\operatorname{Set})]$  spanned by powers of  $\operatorname{id}$ .

Proof. Since  $Z(\mathsf{K})(m,n) = Z(\mathsf{K})(m,1)^n$  and  $[\mathrm{id}^m,\mathrm{id}^n] = [\mathrm{id}^m,\mathrm{id}]^n$ , it is enough to prove the claim in the case n = 1. Let  $f \in Z(\mathsf{K})(m,1)$ . For  $X \in \mathsf{K}(\mathrm{Set})$ , let  $\alpha_{f,X} := Xf : X^m \longrightarrow X$ . We check that this is a morphism of models using lemma 2.3: for  $g: n \longrightarrow 1$ 

$$\begin{aligned} \alpha_{f,X} \circ (X^m g) &= (Xf)(X^m g) & \text{by definition} \\ &= (Xf)(X(g^m \tau_{m;n})) & \text{by lemma 2.17} \\ &= X(fg^m \tau_{m;n}) \\ &= X(gf^n) & \text{since } f \boxtimes g \\ &= (Xg)(Xf)^n \\ &= (Xg)\alpha_{f,X}^n. \end{aligned}$$

For  $F: X \longrightarrow Y$  a morphism of models, we have  $F\alpha_{f,X} = F(Xf) = (Yf)F^n$ 

$$_{f,X} = F(Xf) = (Yf)F^m = \alpha_{f,Y}F^m,$$
 (3.41)

or in other words, the following diagram commutes:

Consequently, the  $\alpha_{f,X}$ 's assemble into a natural transformation  $\alpha_f : \mathrm{id}^m \longrightarrow \mathrm{id}$ . Further, the mapping  $f \longmapsto \alpha_f$  is injective since  $\alpha_{f,K_m} : K_m^m \longrightarrow K_m$  maps  $\mathrm{id}_m$  to f.

Conversely, let  $\beta : \mathrm{id}^m \longrightarrow \mathrm{id}$  be a natural transformation. Let  $b \coloneqq \beta_{K_m}(\mathrm{id}_m) : m \longrightarrow 1$ . We show that for  $X \in \mathsf{K}(\mathrm{Set})$ , the morphism  $\beta_X$  is in fact  $Xb : X^m \longrightarrow X$ . Let  $x \in X^m$ , and  $\tilde{x} : K_m \longrightarrow X$  be the corresponding morphism under the Yoneda lemma. Then by naturality of  $\beta$ , the following square commutes:

$$\begin{array}{cccc}
K_m^m & \xrightarrow{\beta_{K_m}} & K_m \\
\tilde{x}^m & & & \downarrow \tilde{x} \\
X^m & \xrightarrow{\beta_X} & X.
\end{array}$$
(3.43)

Thus,

$$\beta_X(x) = \beta_X \tilde{x}^m(\mathrm{id}_m) = \tilde{x}\beta_{K_m}(\mathrm{id}_m) = \tilde{x}(a) = (Xa)(x).$$
(3.44)

Consequently,  $\beta = \alpha_b$ , and the map  $\alpha : \mathsf{K}(m, 1) \longrightarrow [\mathrm{id}^m, \mathrm{id}]$  is a bijection.  $\Box$ 

# 4. Morita theory

Let L be a Lawvere theory. For  $k \in \mathbb{N}$ , the matrix theory [ARV11, definition 15.4]  $\mathsf{L}^{[k]}$  is the Lawvere theory where  $\mathsf{L}^{[k]}(m,n) = \mathsf{L}(km,kn)$ . The structural morphism  $\aleph_0^{\mathrm{op}} \longrightarrow \mathsf{L}^{[k]}$  maps the projection  $\pi_i : n \longrightarrow 1$  to  $\langle \pi_{k(i-1)+1}, \pi_{k(i-1)+2}, \ldots, \pi_{ki} \rangle : kn \longrightarrow k$ .

Let now  $u: 1 \longrightarrow 1$  be an idempotent operation of L. It is *pseudoinvertible* if there exists  $k \in \mathbb{N}$ ,  $a: 1 \longrightarrow k$  and  $b: k \longrightarrow 1$  such that

$$bu^{\kappa}a = \mathrm{id}_1. \tag{4.1}$$

The *idempotent modification* uLu of L is the subcategory of L spanned by those morphisms  $f: m \longrightarrow n$  such that  $fu^m = f = u^n f$ . The identity of m in uLu is  $u^m$ , and the *i*-th projection  $m \longrightarrow 1$  is  $u\pi_i$ .

**Lemma 4.2.** The idempotent modification uLu is generated by morphisms of the form  $ufu^n$ , where  $f: n \longrightarrow 1$  and  $n \in \mathbb{N}$ .

*Proof.* If  $g: m \longrightarrow n$  is in uLu, then

$$g = u^n g u^m = u^n \langle \pi_i g \rangle_{1 \le i \le n} u^m = \langle u(\pi_i g) u^m \rangle_{1 \le i \le n}.$$

Two Lawvere theories  $K, L \in \mathcal{L}$ aw are *Morita equivalent* [ARV11, definition 15.2], denoted by  $K \sim L$ , if their categories of models K(Set) and L(Set) are equivalent.

**Lemma 4.3.** Let  $K, K', L, L' \in \mathcal{L}$ aw. If  $K \sim K'$  and  $L \sim L'$ , then  $K \otimes L \sim K' \otimes L'$ .

*Proof.* Follows from theorem 3.10 and corollary 3.12.

**Proposition 4.4.** Two commutative  $K, L \in \mathcal{L}$ aw are Morita equivalent if and only if they are isomorphic.

*Proof.* Surely, if K and L are isomorphic, then they are Morita equivalent. Conversely, by proposition 3.40,

$$\mathsf{K}(m,n) = Z(\mathsf{K})(m,n) = [\mathrm{id}_{\mathsf{K}(\mathrm{Set})}^m, \mathrm{id}_{\mathsf{K}(\mathrm{Set})}^n] \cong [\mathrm{id}_{\mathsf{L}(\mathrm{Set})}^m, \mathrm{id}_{\mathsf{L}(\mathrm{Set})}^n] = \mathsf{L}(m,n).$$
(4.5)

**Theorem 4.6** ([ARV11, theorem 15.7]). Let  $\mathsf{K}, \mathsf{L} \in \mathcal{L}$ aw. We have  $\mathsf{K} \sim \mathsf{L}$  if and only if  $\mathsf{K} \simeq u\mathsf{L}^{[k]}u$ , for some  $k \ge 1$  and u a pseudoinvertible idempotent of the matrix theory  $\mathsf{L}^{[k]}$ .

### 5. Stability

A Lawvere theory K is said to be *syntactically stable* at rank  $k \in \mathbb{N}$  if the canonical map  $i_{\mathsf{K}^{\otimes k}} : \mathsf{K}^{\otimes k} \longrightarrow \mathsf{K}^{\otimes k} \otimes \mathsf{K} = \mathsf{K}^{\otimes k+1}$  is an equivalence of categories. Similarly, K is *semantically stable* at rank k if the forgetful functor  $i_{\mathsf{K}^{\otimes k}}^* : \mathsf{K}^{\otimes k+1}(\mathsf{Set}) \longrightarrow \mathsf{K}^{\otimes k}(\mathsf{Set})$  is an equivalence of categories. Of course, if K is syntactically stable at rank k, then it is semantically stable as well. We say that K is syntactically or semantically stable if it is so at rank 1.

**Proposition 5.1.** If K is commutative and semantically stable, then it is syntactically stable.

*Proof.* Follows from lemma 3.39 and proposition 4.4.  $\Box$ 

**Proposition 5.2.** If K is syntactically stable, then it is commutative.

*Proof.* Denote by  $i_1, i_2 : \mathsf{K} \longrightarrow \mathsf{K} \otimes \mathsf{K}$  the canonical morphisms into the left and right component, respectively, and let  $\sigma : \mathsf{K} \otimes \mathsf{K} \longrightarrow \mathsf{K} \otimes \mathsf{K}$  be the symmetry involution, i.e. such that  $\sigma i_1 = i_2$  and conversely. Then by assumption,  $i_1$  is an isomorphisms, and thus, so is  $i_2$ . By definition, for  $f, g \in \mathsf{K}/1$ , we have  $i_1(f) \boxtimes i_2(g)$ , and applying  $i_1^{-1}$  gives  $f \boxtimes i_1^{-1}i_2(g)$ . Note that  $i_1^{-1}i_2 = i_1^{-1}\sigma i_1 = i_2^{-1}i_1$  is an involution. In particular,

$$f \boxtimes i_1^{-1} i_2(i_1^{-1} i_2(g)) = g, \tag{5.3}$$

for all  $f, g \in \mathsf{K}/1$ .

Unfortunately, the converse of proposition 5.2 does not hold, see e.g.  $Mag_1$ .

### 6. Conclusion

TODO: Write the conclusion.

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